Alphabets, Strings, and Languages

Chapter 2

Alphabets and Strings

An *alphabet* Σ is a finite set of symbols.

A *string* or *word* is a finite sequence, possibly empty, of symbols drawn from some alphabet Σ .

- ϵ is the empty string.
- Σ^* is the set of all possible strings over the alphabet Σ .

Alphabet name	Alphabet symbols	Example strings
English alphabet	{a, b, c,, z}	ϵ , aabbcg, aaaaa
Binary alphabet	{0, 1}	ε, 0, 001100
A star alphabet	{★, ❹, ☆, ☆, ☆, ☆}	ε, ΟΟ, Ο★☆☆★☆

String Operations

Length:|s| is the number of symbols in string s. $|\varepsilon| = 0$ |1001101| = 7

Concatenation: *xy* is the **concatenation** of *x* and *y*. If x = good and y = bye, then xy = goodbye ε is the identity for concatenation. So, $x \varepsilon = \varepsilon x = x$

Replication: For each string *w* and each natural number i, the string w^i is defined as follows: $w^0 = \varepsilon$

And if i > 0:

2. 0. 0

 $w^{i} = i$ copies of w concatenated together

So:
$$a^3 = aaa$$

(bye)² = byebye
 $a^0b^3 = bbb$

More String Operations

Reverse: For each string w, w^{R} is the reverse of w.

Concatenation and Reverse: If *w* and *x* are strings, then $(wx)^{R} = x^{R} w^{R}$.

Example: $(nametag)^{R} = (tag)^{R} (name)^{R} = gateman$

Substrings

The string *x* is a **substring** of string *w* if it is a contiguous sequence of characters in *w*. If |x| < |w|, it is a *proper substring*.

Note 1: Every string is a substring of itself.

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Note 2: The empty string ϵ is a substring of every string.

Prefixes

p is a *prefix* of *t* if *t* = *px*

p is a *proper prefix* of *t* iff: *p* is a prefix of *t* and $p \neq t$.

Note 1: Every string is a prefix of itself.

Note 2: The empty string ϵ is a prefix of every string.

Examples:

The prefixes of abbb are: ϵ , a, ab, abb, abbb.The proper prefixes of abbb are: ϵ , a, ab, abb.

Suffixes

s is a **suffix** of *t* if *t* = *x*s

s is a **proper suffix** of t iff: s is a suffix of t and $s \neq t$.

Note 1: Every string is a suffix of itself.

Note 2: The empty string ϵ is a suffix of every string.

Examples:

The suffixes of abbb are: ϵ , b, bb, bbb, abbb.The proper suffixes of abbb are: ϵ , b, bb, bbb.

Sets

- The set with no members is the *empty set*, denoted {} or \emptyset
- If every element of set A is a member of set B, we say that A is a subset of B, denoted $A \subseteq B$
- If set A is a subset of set B and A ≠ B, then A is a proper subset of B.
- The empty set is a subset of any set.
- Set membership: $x \in S$ denotes that x is in the set **S**

More on Sets

- The *union* of two sets, $A \cup B$, is the set that contains everything that is in *A*, in *B*, or in both.
- The *intersection* of two sets, $A \cap B$, is the set that contains only those elements that are in both A and B.
- Two sets A and B are **disjoint** if they have no elements in common, i.e. $A \cap B = \emptyset$
- The set difference of A and B, A B, is the set that contains everything that is in A but not in B.

Even More on Sets

- The complement of A, written A, is the set containing everything that is in some universal set U that is not in A.
- The *cardinality* of a set A, written |A|, is the number of elements in A.
- The *power set* of *A*, *denoted P(A)*, is the set of all subsets of *A*. Note that |P(A)| = 2^{|A|}.

Defining a Language

A *language* is a (finite or infinite) set of strings over an alphabet Σ .

Examples:

Let $\Sigma = \{a, b\}$

Some languages over Σ :

 \emptyset , (the language with no strings in it) { ϵ }, (the language containing just the empty string) {a, b}, { ϵ , a, aa, aaa, aaaa, aaaaa}



How Large is a Language?

The smallest language over any Σ is \emptyset , i.e. the language with no strings in it, so size 0.

The language Σ^* contains a countably infinite number of strings: the empty string ε and all strings of any length containing only the characters a and b.

Example Language Definition

 $L = \{x \in \{a, b\}^* : all a's precede all b's\}$

Then:

abbb, aabb, and aaaaaaabbbbb are in L.

aba, ba, and abc are not in L.

What about: ϵ , a, aa, and bb?

Exercises

What are the following languages?

$$L = \{x : \exists y \in \{a, b\}^* : x = ya\}$$

 $L = \{w \in \{a, b\}^*: no \text{ prefix of } w \text{ contains } b\}$

 $L = \{w \in \{a, b\}^*: no \text{ prefix of } w \text{ starts with } a\}$

 $L = \{w \in \{a, b\}^*: every prefix of w starts with a\}$

Functions on Languages

- Set operations
 - Union
 - Intersection
 - Complement
- Language operations
 - Concatenation
 - Kleene star

Concatenation of Languages

If L_1 and L_2 are languages over Σ :

 $L_1L_2 = \{ w \in \Sigma^* : \exists s \in L_1 \text{ and } \exists t \in L_2 \text{ such that } w = st \}$

Example: $L_1 = \{cat, dog\}$ $L_2 = \{apple, pear\}$ $L_1 L_2 = \{catapple, catpear, dogapple, dogpear\}$

{ ϵ } is the identity for concatenation: $L{\epsilon} = {\epsilon}L = L$

 \varnothing is the "zero" for concatenation: $L \oslash = \oslash L = \oslash$

Be Careful Concatenating Languages Defined Using Variables

 $L_1 = \{a^n: n \ge 0\} \\ L_2 = \{b^n : n \ge 0\}$

 $L_1 L_2 = \{a^n b^m : n, m \ge 0\}$ $L_1 L_2 \neq \{a^n b^n : n \ge 0\}$

Kleene Star

$$L^* = \{\varepsilon\} \cup \\ \{w \in \Sigma^* : w = w_1 \ w_2 \ \dots \ w_k \text{ and } \\ w_1, \ w_2, \ \dots \ w_k \in L\}$$

Example: $L = \{ dog, cat, fish \}$ $L^* = \{ \epsilon, dog, cat, fish, dogdog, dogcat, fishcatfish, fishdogdogfishcat, ... \}$

 $L^+ = L L^*$

 $L^+ = L^* - \{\varepsilon\}$ iff $\varepsilon \notin L$

Lexicographic Ordering

Lexicographic Order:

- Shortest first
- Within a given length, "dictionary" order

What is the lexicographic enumeration of:

 $\{w \in \{a, b\}^* : |w| \text{ is even}\}$